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Heterogeneous Concurrent Error Detection (*hCED*) Based on Output Anticipation

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Agenda



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- **Related Work**
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- **The Baseline Arrangement**
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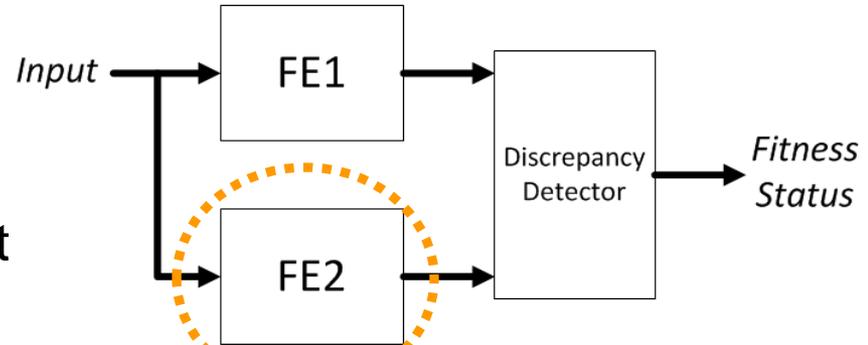


Introduction

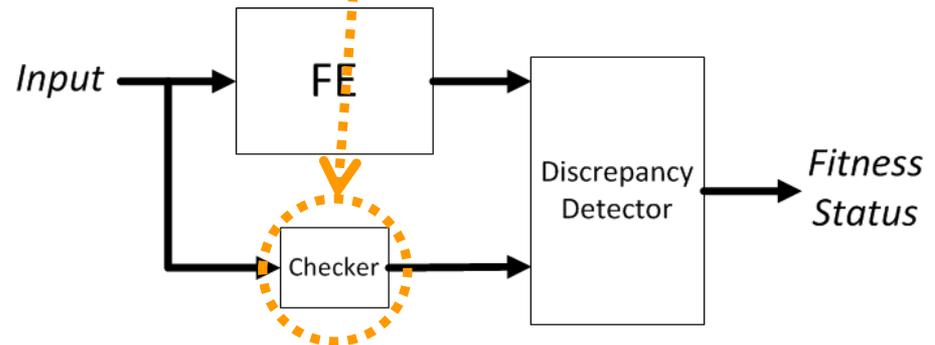


- Redundancy-based **Fault Detection** → mission critical applications
- A conventional **Concurrent Error Detection (CED)** technique → Two **identical** replicas of a given module are required to provide fault detection capability [1],[2]
- A **Triple Modular Redundancy (TMR)** based design → faults can be handled faults via majority **Voting** [3]
- The redundancy based schemes have **overhead** of resource and power [4]

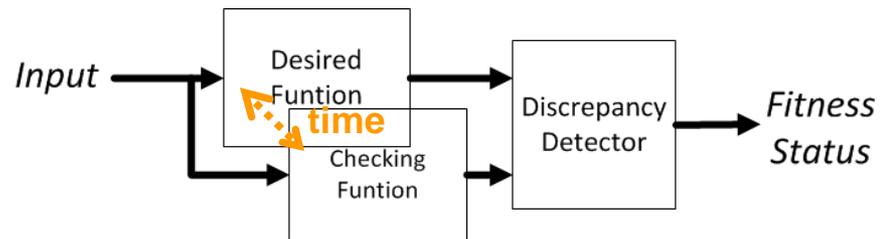
Conventional Concurrent
Error Detection arrangement



Heterogenous Spatial
Concurrent Error
Detection arrangement



Heterogenous Temporal
Concurrent Error
Detection arrangement





Related Work



- A research of different **forms** of CED arrangement was presented by Mitra *et. al* [1] → e.g. **Parity-Based** systems
- Built-in Self Test (BIST) → **Detection Latency**
- *Temporal* redundancy techniques [11] → A typical error detection scheme involves running a duplicate thread for the comparison purpose on a **chip multiprocessor (CMP)**
 - Effective for transient errors, but **permanents faults** are not detectable
- A **resource testing** scheme by Gao *et. al* [15] using time multiplexing of different components through the **reconfiguration capability** of FPGA → drawbacks of BIST along with the **reconfiguration times**



Salient Features of the Proposed Work



- As the **reconfiguration time** is a considerable entity in current FPGA technology, we propose to **multiplex** the inputs to a fixed hardware fabric instead of reconfiguring the resources with alternating functions.

Spatial *h*CED:

- Fault detection in Spatial *h*CED mode with **resource saving**
→ Area of Checker << Area of Functional Module

Temporal *h*CED:

- Fault detection in Temporal *h*CED mode with **uniplex chip area** requirement at the cost of reduced throughput.
→ Same fabric resources are time-multiplexed across different inputs
- The coverage of transients as well as **permanent faults** in the fault detection using Temporal *h*CED.
→ Since Checker design differs from Functional Module design, based on the inputs



Case Study - DCT hardware Core



- Common in **image/video compression applications**, and hardware implementation is highly desired in many applications due to parallel nature of the image processing related tasks and their **throughput requirements**
- **Video encoder application** → image frame is first divided into Macroblocks and the **transform** operation is performed on these Macroblocks
- **DCT Transformation** uses following kernels [16],[17]:

$$r(x, y, u, v) = \alpha(u)\alpha(v)\cos\frac{(2x+1)u\pi}{2n}\cos\frac{(2y+1)v\pi}{2n} \quad (2)$$

where

$$\alpha(u) = \begin{pmatrix} \sqrt{\frac{1}{n}} & \text{for } u = 0 \\ \sqrt{\frac{2}{n}} & \text{for } u = 1, 2, \dots, n-1 \end{pmatrix}$$



A Necessary Condition for Fault-free DCT



The **1D DCT operation** on a row of pixels in a Macroblock is defined by:

$$Y = \Phi.X$$

Where

X = A row of input pixels macroblock.i.e., $\{x_1, x_2, \dots, x_8\}$

Φ = Matrix of DCT kernels

Y = A row of output DCT coefficients.i.e., $\{y_1, y_2, \dots, y_8\}$

As the DCT operation is a **linear transformation** from input pixels space to output coefficients space, we can make the following derivations:

$$(Y - \hat{Y}) = \Phi.(X - \hat{X}) \quad \Delta X = X - \hat{X}, \Delta Y = Y - \hat{Y}$$


$$\Delta Y = \Phi.\Delta X$$

The current DCT coefficients can be **estimated** from the previous coefficients using the difference values. $\hat{Y} = \hat{Y} + \Delta Y$

A **necessary condition** for a fault free DCT block can be written as:

$$Y = \hat{Y}$$

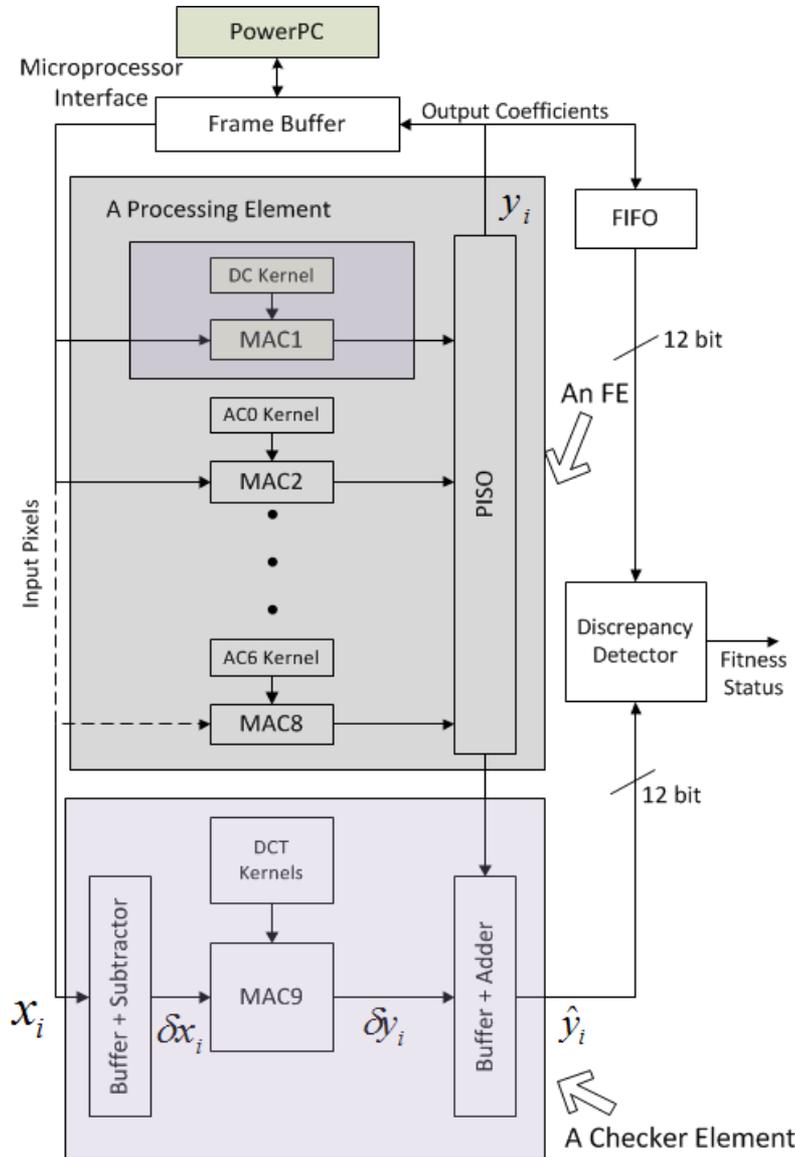
Estimated = Actual



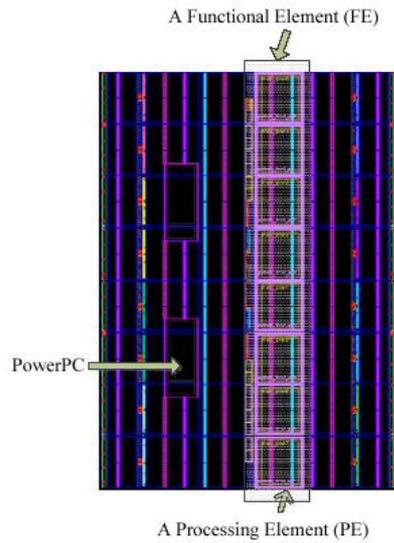
Baseline Arrangement



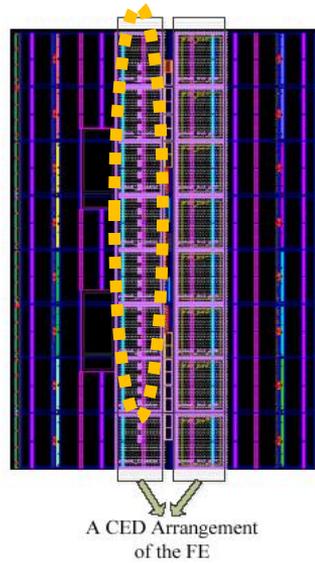
- An **8-point 1-D DCT** implemented in Verilog HDL using Xilinx ISE 9.2i development environment.
- Xilinx PlanAhead for the partial reconfiguration design
- Xilinx development board ML410 , Virtex-4 FPGA.
- Multiple **Processing Elements (PEs)** used to compute DCT coefficients, **one coefficient per PE**
- A **Functional Element (FE)** consists of 8 *Processing Elements* (PEs), where a FE will compute the 8-point 1D DCT.
- For example, PE1 contains the **DC-kernel**, PE2 has the **AC0-kernel** and so on.
- The DCT core is interfaced with on-chip PowerPC microprocessor through GPIO core.



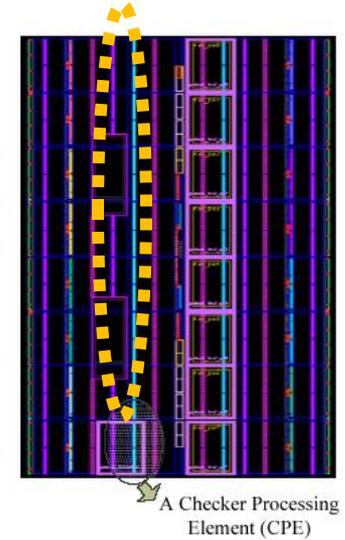
Compared to conventional CED Arrangement, the Checking Element contains a **single PE** instead of 8 PEs.



(a) An FE realizing the DCT function



(b) CED arrangement



(c) Spatial *h*CED arrangement



Spatial Heterogeneous CED



- PE1 through PE8 form an FE computing the DCT of input data.
- The PE containing a MAC9 serves as a **Checker Processing Element (CPE)**. The CPE performs MAC operation on the difference input, instead of the input pixel values.
- The CPE utilizes the previous output from FE to predict the current output → discrepancy reveals error

RESOURCE UTILIZATION FOR SPATIAL HCED ARRANGEMENT

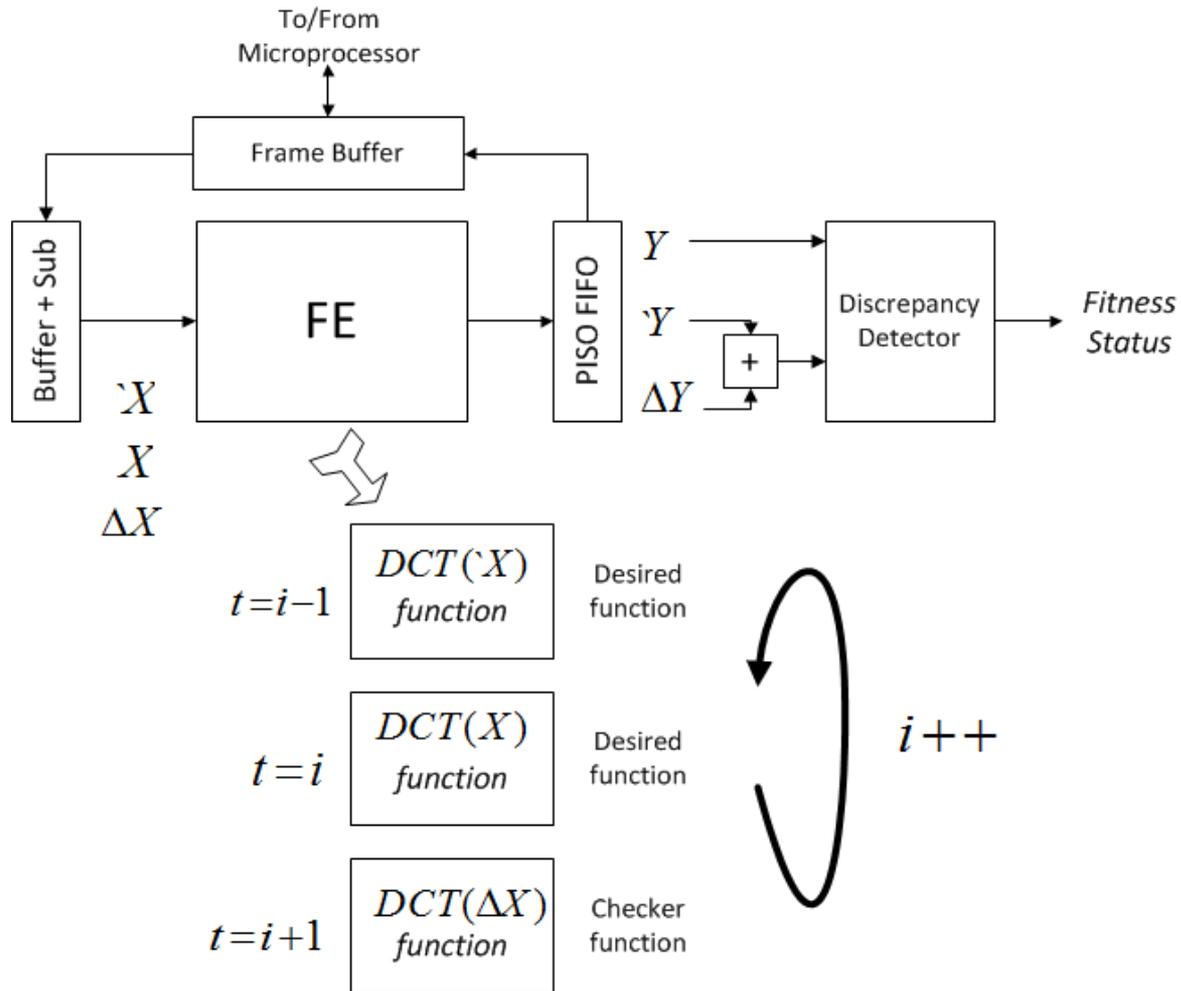
Component	PE	FE	CPE
Number of slices	91	728	167
Number of slice FFs	46	368	75
Number of 4 input LUTs	161	1288	308



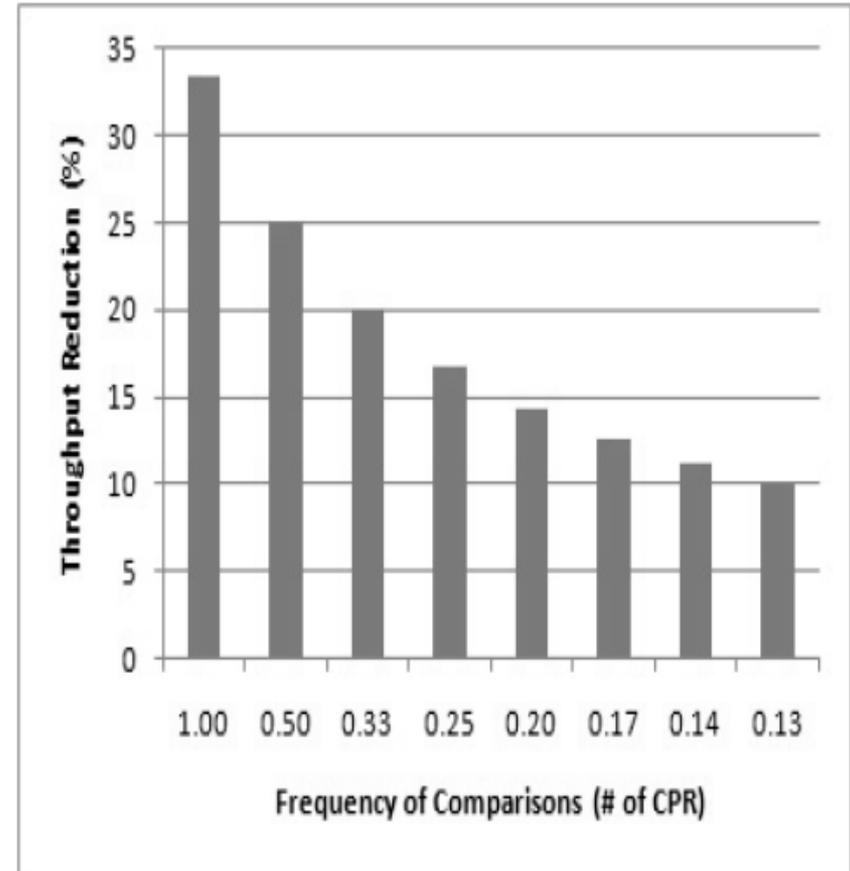
Spatial Heterogeneous CED



- The 8 PEs utilize 728 Configuration Logic Block (CLB) slices of the chip, while one CPE implementation requires only 167 slices, saving **overall 77%** of the FPGA resource.
- A **conventional CED** would require approximately **2X resource count**
- One PE requires an estimated power of 12.39 mW. On the other hand, a CPE's estimated dynamic power consumption is 13.33 mW using a clock frequency of 100 MHz → **Power saving 1/8th**
- Inject **Stuck-At faults** at LUT inputs and analyze their behavior in post Place-and-Route simulation model.
- Saving in resource and power at the expense of fault detection latency compared to conventional CED



- **Throughput reduction** of *Temporal hCED* arrangement
- For example, if one time instant out of 3 instants is dedicated for the prediction computation, the *effective throughput* would **reduce to 66.7%**
- By performing the redundant computation of difference input data **less frequently**, the useful throughput can be improved from the worst case.





Conclusions



- *hCED* provides a **flexible** fault detection scheme.
- The ***spatial hCED*** form exhibits reduced resource requirements than the conventional CED technique. Thus, **area and power** can be saved using the proposed approach at the cost of a negligible fault detection latency overhead.
- The ***temporal hCED*** form has error detection capability of fault coverage that includes **permanent faults** in logic resources, in addition to transient faults.
- The ***temporal*** error detection form has **uniplex area** requirement avoiding redundancy in the resources at the cost of throughput reduction.
- The *hCED* scheme appears to be readily applicable to any **linear transformation** that is pipelineable.



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